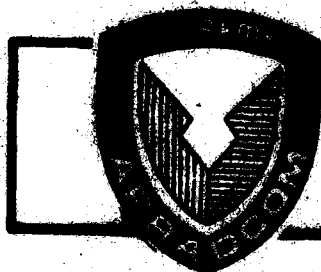TECHNICAL REPORT ARLCB-TR-82028A

SALOME, A STRUCTURED AND LOGICALLY MINIMAL ENSEMBLE OF PROGRAMMING CONSTRUCTS

Royce W. Soanes, Jr.

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US ARMY ARMAMENT RESEARCH AND DEVELOPMENT COMMAND
LARGE CALIBER WEAPON SYSTEMS LABORATORY
BENET WEAPONS LABORATORY
WATERVLIET, N. Y. 12189

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19. KEY WORDS (Continue on reverse side if necessary and identify by block number)

Structured Language	Strings	Program Generator
Structured Programming	Compiler	Source Language
Language Translation	Precompiler	Target Language
Language Translator		

20. ABSTRACT (Continue on reverse side if necessary and identify by block number)

An overview of the Salome structured programming language will be given, along with an appendix briefly describing its syntax and semantics. Salome has been designed with the Fortran programmer in mind. The Salome source language is supported by a translator having a target language of standard Fortran, and Fortran code may be injected into Salome source code when needed. The Salome translator has been written in Salome and Fortran.

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OVERVIEW

The acronym SALOME stands for structured and logically minimal ensemble (or selection and looping operations made easy). The term "structured" refers to the gotoless nature of the logical control statements of Salome. While the power of the goto statement has corrupted many programs, the logical control statements of Salome virtually eliminate the need to use goto statements and labels in programs. The phrase "logically minimal" refers to the fact that the variety and semantic ambiguity of the structured logical control statements are kept to a minimum.

There is one looping construct and one selection construct in Salome, while other languages may force the programmer to learn several. Flexibility and unambiguous semantics are stressed in Salome - not variety.

Blanks are as important in Salome as they are in ordinary English text. The use of blanks as delimiters enhances readability, eliminates extraneous punctuation, and allows for greater brevity of syntax. Although Salome is delimiter oriented (as opposed to line oriented) there is no general end of statement or between statement delimiter in Salome. A small penalty one pays for this feature is that blanks are not allowed in assignment statements. This was considered to be a small price to pay for the elimination of a lot of extraneous semicolons.

The Salome language is supported by a one pass translator whose target language is Fortran; hence, when one writes a program in Salome, one is in effect writing in two different high level languages at the same time. The translator is in turn supported by a package of Fortran callable string manipulation routines. The primary reason Fortran was selected as the target

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language was to benefit those unfortunate programmers who are inextricably mired down in a Fortran environment when Fortran is more lacking in structured programming constructs than any other high level language except Basic. A widely used high level target language also assures greater portability of Salome and its translator.

The file produced by the Salome translator consists of the original Salome source code inserted as special Fortran comments interleaved with the generated Fortran code. When Salome syntax errors are caught, they are pointed out by a string of dots. If there are syntax errors in a Salome routine, only the first one in that routine is flagged and the rest of the code in that routine is ignored. No attempt at error recovery is made because one can't be absolutely positive about what the programmer actually intended, and guessing usually uncovers syntax errors that aren't.

The reason for the interleaving of the Salome source code as comments with the generated Fortran code is to enable the programmer to relate any syntax errors picked up by the Fortran compiler (and not picked up by the Salome translator) back to the original Salome source code. The Salome translator does not check for things that the Fortran compiler is going to check for anyway.

Since the appendix of this paper contains a brief but fairly complete description of Salome, no further details will be described here. Instead, a small subroutine written in Salome along with the generated Fortran and the interleaved file are presented as a small exercise in deduction. If one knows Fortran and one reads the interleaved file carefully, one can deduce the exact semantic meaning of the Salome routine.

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SALOME SOURCE EXAMPLE

```
SUB LININT ( N X Y XI YI ) -- --------------------------------------LININT
-- LINEAR INTERPOLATION
(--
N=NO. OF POINTS.
X=ABSCISSA ARRAY.
Y=ORDINATE ARRAY.
XI=ABSCISSA AT WHICH INTERPOLATION IS DESIRED.
YI=INTERPOLATED RESULT.
X IS ASSUMED TO BE SORTED IN ASCENDING ORDER.
--)
DIM X 1 , Y 1 .
-- ASSUME AT LEAST 2 POINTS OF DATA
@ N > 1 @
-- GET PROPER SUBINTERVAL USING BINARY SEARCH
IL=1 IR=N
DO # IL+1 = IR #
    -- COMPUTE INDEX OF ABSCISSA MIDWAY BETWEEN IL AND IR
    IM=(IL+IR)/2
    -- REDEFINE IL OR IR AS IM
    IF XI < X(IM) , IR=IM ; IL=IM FI
    -- ASSUME IL AND IR ARE STILL IN PROPER ORDER
    @ IL < IR @ OD
DX=X(IR)-X(IL) DY=Y(IR)-Y(IL)
-- ASSUME LENGTH OF SUBINTERVAL IS POSITIVE
@ DX > 0. @
YI=Y(IL)+(DY/DX)*(XI-X(IL))
RET END
```

FORTRAN GENERATED FROM SALOME SOURCE EXAMPLE

```fortran
      SUBROUTINE LININT(N,X,Y,XI,YI)
C     ---------------------------------------------LININT
C     LINEAR INTERPOLATION
C
C     N=NO. OF POINTS.
C     X=ABSCISSA ARRAY.
C     Y=ORDINATE ARRAY.
C     XI=ABSCISSA AT WHICH INTERPOLATION IS DESIRED.
C     YI=INTERPOLATED RESULT.
C     X IS ASSUMED TO BE SORTED IN ASCENDING ORDER.
C
      DIMENSION X(1),Y(1)
C     ASSUME AT LEAST 2 POINTS OF DATA
      IF(N.GT.1) GO TO 1000
      WRITE(6,1001)
 1001 FORMAT(/////,' N > 1 IS FALSE IN LININT')
      CALL EXIT
C     GET PROPER SUBINTERVAL USING BINARY SEARCH
 1000 IL=1
      IR=N
 1002 IF(IL+1.EQ.IR) GO TO 1003
C     COMPUTE INDEX OF ABSCISSA MIDWAY BETWEEN IL AND IR
      IM=(IL+IR)/2
C     REDEFINE IL OR IR AS IM
      IF(.NOT.(XI.LT.X(IM))) GO TO 1005
      IR=IM
      GO TO 1004
 1005 IL=IM
C     ASSUME IL AND IR ARE STILL IN PROPER ORDER
 1004 IF(IL.LT.IR) GO TO 1006
      WRITE(6,1007)
 1007 FORMAT(/////,' IL < IR IS FALSE IN LININT')
      CALL EXIT
 1006 GO TO 1002
 1003 DX=X(IR)-X(IL)
      DY=Y(IR)-Y(IL)
C     ASSUME LENGTH OF SUBINTERVAL IS POSITIVE
      IF(DX.GT.0) GO TO 1008
      WRITE(6,1009)
 1009 FORMAT(/////,' DX > 0.  IS FALSE IN LININT')
      CALL EXIT
 1008 YI=Y(IL)+(DY/DX)*(XI-X(IL))
      RETURN
      END
```

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```
C__S__SUB LININT ( N X Y XI YI ) — ---------------------------------LININT
      SUBROUTINE LININT (N,X,Y,XI,YI)
C       ------------------------------------------LININT
C__S__-- LINEAR INTERPOLATION
C     LINEAR INTERPOLATION
C__S__(--
C
C__S__N=NO. OF POINTS.
C     N=NO. OF POINTS.
C__S__X=ABSCISSA ARRAY.
C     X=ABSCISSA ARRAY.
C__S__Y=ORDINATE ARRAY.
C     Y=ORDINATE ARRAY.
C__S__XI=ABSCISSA AT 'WHICH INTERPOLATION IS DESIRED.
C     XI=ABSCISSA AT WHICH INTERPOLATION IS DESIRED.
C__S__YI=INTERPOLATED RESULT.
C     YI=INTERPOLATED RESULT.
C__S__X IS ASSUMED TO BE SORTED IN ASCENDING ORDER.
C     X IS ASSUMED TO BE SORTED IN ASCENDING ORDER.
C__S__--)
C
C__S__DIM X 1 , Y 1 .
      DIMENSION X(1),Y(1)
C__S__-- ASSUME AT LEAST 2 POINTS OF DATA
C     ASSUME AT LEAST 2 POINTS OF DATA
C__S__@ N > 1 @
      IF(N.GT.1) GO TO 1000
      WRITE(6,1001)
 1001 FORMAT(/////,' N > 1 IS FALSE IN LININT')
      CALL EXIT
C__S__-- GET PROPER SUBINTERVAL USING BINARY SEARCH
C     GET PROPER SUBINTERVAL USING BINARY SEARCH
C__S__IL=1 IR=N
 1000 IL=1
      IR=N
C__S__DO # IL+1 = IR #
 1002 IF(IL+1.EQ.IR) GO TO 1003
C__S__   -- COMPUTE INDEX OF ABSCISSA MIDWAY BETWEEN IL AND IR
C     COMPUTE INDEX OF ABSCISSA MIDWAY BETWEEN IL AND IR
C__S__   IM=(IL+IR)/2
      IM=(IL+IR)/2
C__S__   -- REDEFINE IL OR IR AS IM
C     REDEFINE IL OR IR AS IM
```

```
C__S__    IF XI < X(IM) , IR=IM ; IL=IM FI
      IF(.NOT.(XI.LT.X(IM))) GO TO 1005
      IR=IM
      GO TO 1004
 1005 IL=IM
C__S__    -- ASSUME IL AND IR ARE STILL IN PROPER ORDER
C     ASSUME IL AND IR ARE STILL IN PROPER ORDER
C__S__    @ IL < IR @ OD
 1004 IF(IL.LT.IR) GO TO 1006
      WRITE(6,1007)
 1007 FORMAT(/////,' IL < IR IS FALSE IN LININT')
      CALL EXIT
 1006 GO TO 1002
C__S__DX=X(IR)-X(IL) DY=Y(IR)-Y(IL)
 1003 DX=X(IR)-X(IL)
      DY=Y(IR)-Y(IL)
C__S__-- ASSUME LENGTH OF SUBINTERVAL IS POSITIVE
C     ASSUME LENGTH OF SUBINTERVAL IS POSITIVE
C__S__@ DX > 0. @
      IF(DX.GT.0) GO TO 1008
      WRITE(6,1009)
 1009 FORMAT(/////,' DX > 0. IS FALSE IN LININT')
      CALL EXIT
C__S__YI=Y(IL)+(DY/DX)*(XI-X(IL))
 1008 YI=Y(IL)+(DY/DX)*(XI-X(IL))
C__S__RET END
      RETURN
      END
```

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APPENDIX

A QUICK GUIDE TO THE SALOME PROGRAMMING LANGUAGE

1. Definition of Isolated String

    An isolated string resides on a single line, is preceded by one or more

    blanks (or begins in Column 1) and is followed by one or more blanks (or

    ends in Column 72).

    All Salome keywords must be isolated strings containing no blanks.

2. Comments

    A. A 'tack on' comment may stand alone on a line or it may be tacked ont

        the end of another statement.

        Ex.

        -- TACK ON COMMENT STANDING ALONE
        A=B+C -- TACKED ON TACK ON COMMENT

        '--' is the keyword for tack on comments

        The tack on comment has essentially the same form as the ADA comment.

    B. A delimited comment may occupy more than one line.

        Ex.

        (-- THIS IS
        A DELIMITED
        COMMENT--)

        '(--' AND '--)' are the keywords for delimited comments.

        Delimited comments may be used to temporarily disable certain sections

        of executable code and they may be nested.

3. Injected Fortran

Fortran statements may be injected into a Salome program when needed.

Ex.

```
F       COMMON A(100),B(50)
F       A = B * C
```

'F' is the keyword for injecting Fortran. The Fortran statement

begins immediately after the blank following 'F'.

4. Variable Declaration Statements

The correspondence between Salome variable declaration keywords and

Fortran variable declaration keywords is given by the following table.

SALOME	FORTRAN
DIM	DIMENSION
INT	INTEGER
REAL	REAL
DP	DOUBLE PRECISION
LOG	LOGICAL
EQV	EQUIVALENCE

Ex.

```
DIM X 10 20 , Y 100 . translates to
DIMENSION X(10,20),Y(100)

INT A , B , C 10 20 30 . translates to
INTEGER A,B,C(10,20,30)
```

The aforementioned keywords plus ',' and '.' are the keywords for

declaration statements.

'.' Ends all variable declaration statements.

',' Separates variable references in all variable declaration statements

except the equivalence statement.

',' Separates variable group references in the equivalence statement.

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Ex.

```
EQV A B C , K(5) L M(10) .  Translates to
     EQUIVALENCE (A,B,C),(K(5),L,M(10))
```

Note that a string such as A(10,20) may appear in an equivalence statement

but not in any other variable declaration statement.

5. Assignment Statements

Salome assignment statements must be isolated strings containing no

blanks. They are otherwise equivalent to Fortran assignment statements.

If blanks are deaired in assignment statements, they may be injected as

Fortran.

Ex.

A=B*C+D P=Q*R

F   A = B + C * D / E

6. I/O

A. The keywords which begin I/O statements are:

'IN' 'OUT' 'R' 'W' 'RF' 'WF' 'RU' 'WU' 'FMT'

They may be read as follows:

```
IN - Input on device no.
OUT - Output on device no.
R - Read
W - Write
RF - Read in Format
WF - Write in Format
RU - Read Unformatted
WU - Write Unformatted
FMT - Format
```

B. The rest of the I/O keywords are:

'(' ')' '.' 'EOF:' 'ERR:'

'(' and ')' bracket format strings in 'R' and 'W' statements.

'.' ends all I/O statements except for the IN and OUT statements.

'EOF:' precedes any end of file label name.

'ERR:' precedes any error label name.

Ex.

IN 1

OUT IOUT

R I X Y (I3 2F10.0) EOF: END-OF-DATA .

W I X Y (' I=' I3 ' X=' E14.7 ' Y=' E14.7) .

RF F1 A B EOF: ENDDATA .
FMT F1 8F10.0 .

W (' PLAIN HOLLERITH ') .

Note that format strings in 'FMT' statements are not enclosed in

parentheses as they are in 'R' and 'W' statements.

Format references in RF, WF, and FMT statements may be any isolated

string containing no blanks.

7. External Subroutine Calls

   External subroutines are called by writing the name of the subroutine
   follwed by its argument list, if any. No 'call' keyword is used and
   English Salome keywords may not be used as subroutine names.

   Ex.

   EXIT - CALL TO EXIT

   ADD (A B C) — ADD A TO B GIVING C

   DEFSTR (ABC ' A B C ') — DEFINE STRING

   The only keywords associated with an external subroutine call are '(' and
   ')'.

   Arguments in an external subroutine call must be isolated strings. The
   only argument which may contain blanks is quoted Hollerith. Arguments are
   not separated by commas.

8. External Subprogram Declaration Statements

   The keywords beginning external subprogram declarations and their Fortran
   counterparts are given by the following table.

SALOME	FORTRAN
SUB	SUBROUTINE
FUN	FUNCTION
IFUN	INTEGER FUNCTION
RFUN	REAL FUNCTION
LFUN	LOGICAL FUNCTION
DPFUN	DOUBLE PRECISION FUNCTION

   '(' and ')' are the only other external subprogram declaration keywords.

Additional keywords necessary to complete the definition of an external

subprogram are:

SALOME   FORTRAN

==>   ENTRY
RET   RETURN
END   END

Ex.

SUB INIT — INITIALIZATION ROUTINE

SUB MULT (A B C) — MULTIPLY B BY A GIVING C

 FUN PROD (X Y) — Product of X and Y

Parameters in an external subprogram declaration must be isolated strings

containing no blanks.

9. Internal Subroutine Calls

External subroutine calls in Salome are quite similar to those of Fortran,

but Salome also has the facility for defining and calling internal

subroutines. These sections of code are called internal because they

belong to the program or subprogram which uses them and they may not be

called by any other external subprogram.

An internal subroutine name may be any isolated string.

An internal subroutine is called by enclosing its name in double quotes

(not two single quotes). The double quotes are the only keywords

associated with the internal subroutine call.

Ex.

" READ INPUT DATA "

" INITIALIZE ARRAYS "

No arguments are passed to internal subroutines. All data in the calling
program is available to the internal subroutine.

10. Internal Subroutine Declaration Statements

Internal subroutines are declared at the end of the calling program

(after a return or call to exit) by writing 'TO' followed by the

internal subroutine name enclosed in double quotes, followed by any code,

followed by 'OT'.

Ex.

```
EXIT ( OR RET )
TO " READ INPUT DATA "
   R X Y Z ( 3F10.0 ) . OT
```

Internal subroutines may call other internal subroutines within the same

program or subprogram.

Internal subroutines have exactly one exit point at 'OT'.

The Salome internal subroutine simplifies the simulation of recursion

considerably.

11. If Statement

The 6 keywords involved in the If Statement are:

'IF' ',' '$' '/' ';' 'FI'

The basic functions of these keywords during translation time and run time are:

'IF' Tells the translator that the first Boolean expression is about to begin.

',' Ends a Boolean expression and selects the statement sequence following to be executed if the Boolean is true. This keyword may be read: 'then'.

'$' Ends a Boolean expression and selects the statement sequence following to be executed if the Boolean is false. This keyword may be read: 'Then Don't Select'.

'/' Halts execution in the if statement if the previous statement sequence has been executed and tells the translator that another Boolean expression is about to begin. This keyword may be read: 'Otherwise,If'.

';' Halts execution in the if statement if the previous statement sequence has been executed and selects the statement sequence following to be executed if no other statement sequence has yet been executed. This keyword may be read 'Otherwise!'.

'FI' ends the If Statement.

The Relational and Boolean operators used in Boolean expressions are:

'<' '>' '<=' '>=' '=' '/=' 'AND' 'OR' 'NOT'

Relational operators must not contain blanks, but they need not be isolated. The Boolean operators 'AND', 'OR' and 'NOT' must be isolated and not contain blanks.

The general form of the If Statement is as follows:

```
IF B1 , S1 /
   B2 , S2 /
   B3 , S3 /
      .
      .
   BI , SI /
      .
      .
   BN , SN ; S FI
```

Where B1 through BN are Boolean expressions and S1 through SN and S are sequences of statements.

If Statements may be nested to theoretically any level.

Ex.

```
IF A>0. AND B>0. , X=A*B ; X=0. FI

IF I/=0 , X=0. Y=0. Z=0. FI

IF X < A , F=-1. /
   X > B , F=1. ; F=0. FI

IF I<1 , I=1 /
   I>N , I=N FI
```

12. Assertion or Assumption Statements

One may very quickly insert error checking code into a program via the
assumption or assertion statement.

This statement consists of nothing more than a Boolean expression
delimited on both sides by the keyword '@'.

Ex.

SUB INTERP (N X Y XI YI) — INTERPOLATION ROUTINE
.
.
@ N>1 @
.
.
END

The assumption that N > 1 in subroutine interp will translate into the
following typical Fortran code.

```
          IF(N.GT.1) GO TO 1000
          WRITE(6,1001)
1001 FORMAT(//////,' N>1 IS FALSE IN INTERP')
          CALL EXIT
1000 CONTINUE
```

The Salome assertion statement is considerably less verbose than the
corresponding Fortran error checking code and will therefore make it far
easier for the programmer to give his program the higher order of
intelligence it needs in order to detect bad data or situations and
subsequently notify the programmer in fairly explicit terms.

The write statement generated by the assertion statement always writes to
6.

13. Goto Statements and Labels

   Although they will seldom be needed in a well written Salome program, goto
   statements and labels are available in Salome.

   The Salome keyword '->' corresponds to the keyword "goto" in Fortran.

   In Fortran, labels must be numbers (statement numbers). Fortran labels
   therefore have no mnemonic value. In Salome, however, label names may be
   any isolated string containing no blanks. Salome labels may therefore be
   given considerable mnemonic value.

   The Salome goto statement contains the goto arrow followed by a label
   name. The corresponding label statement or destination of the goto
   consists of the label name enclosed in the special brackets: '<<' and
   '>>'. This convention makes labels stand out well and is used in the ADA
   programming language. The label statement corresponds to the 'CONTINUE'
   statement in Fortran.

   The only keywords associated with goto and label statements are:

   '->' '<<' '>>'

   Ex.

   .
   .
   << DATA-INPUT-SECTION >>
   .
   .
   -> DATA-INPUT-SECTION
   .
   .
   -> END-OF-PROGRAM
   .
   .
   << END-OF-PROGRAM >>
   END

14. Loops

Salome has a single looping statement with auxiliary loop escape
and loop continue statements.

A Salome loop starts with 'DO' and ends with 'OD'. Whatever code lies
between these two keywords is executed over and over again until some form
of loop escape is done.

The sharp sign (#) is the basic symbol used to indicate loop escapes.

To escape from 1,2,3 or 4 surrounding loops if Boolean expression B is
true, one writes:

```
# B #
    or
## B ##
      or
### B ###
        or
#### B ####
```
Respectively

To escape from 1,2,3 or 4 surrounding loops unconditionally, one writes
the loop escape arrows:

```
-#>
 or
--#>
  or
---#>
   or
----#>
```
Respectively

To escape from more than 4 levels of loop nesting, one must use a goto and
label.

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While a loop escape statement breaks off execution of a loop completely, a

loop continue statement breaks off execution of a single pass or iteration

and then continues execution back at the beginning of the loop. The colon

(:) is the basic symbol used to indicate loop continuation.

In Fortran, when one 'GOES TO' the last statement in a 'DO' loop (often a

'CONTINUE' statement), one is doing loop continuation.

To continue 1,2,3 or 4 surrounding loops if Boolean expression B is true,

one writes:

```
: B :
   or
:: B ::
     or
::: B :::
       or
:::: B ::::
Respectively
```

To continue 1,2,3 or 4 surrounding loops unconditionally, one writes the

loop continue arrows:

```
<:-
 or

<:--
 or
<:---
 or
<:----
Respectively
```

To continue beyond 4 levels of loop nesting, one must use a goto and

label.

A-13

The unconditional loop escape and loop continue statements allow the programmer to perform some action just prior to escape or continuation. The loop continuation constructs are seldom needed, but if they are needed, one should be careful to increment any loop index at the beginning of the loop in order to avoid an endless loop.

Ex.

```
-- MULTIPLY M X N MATRIX A TIMES N X P MATRIX B TO
-- GIVE M X P MATRIX C

-- MAKE SOME EXPLICIT ASSUMPTIONS
@ M>0 @ @ N>0 @ @ P>0 @

(-- IF THESE ASSUMPTIONS ARE NOT MADE EXPLICITLY AND ANY ONE OF THEM
TURNED OUT TO BE FALSE, NO ERROR WOULD OCCUR, BUT MATRIX C WOULD NOT BE
DEFINED AND NEITHER THE REST OF THE PROGRAM NOR THE PROGRAMMER WOULD BE
MADE IMMEDIATELY AWARE OF THIS SITUATION--)

I=0
DO I=I+1 # I>M # -- INDEX ROW OF A
   J=0
   DO J=J+1 # J>P # -- INDEX COLUMN OF B
      K=0 C(I,J)=0.
      DO K=K+1 # K>N # -- INDEX SUMMAND OF ROW A/COL B
                       -- INNER PRODUCT
         C(I,J)=C(I,J)+A(I,K)*B(K,J)
            OD OD OD -- END OF MATRIX MULTIPLICATION

-- LET ARRAY A HAVE N ELEMENTS
-- ADD UP THE POSITIVE ELEMENTS OF ARRAY A
I=0 S=0.
DO I=I+1 # I>N # : A(I)<=0. : S=S+A(I) OD

-- ADD UP THE ELEMENTS OF ARRAY A HAVING INDICES BETWEEN
-- OR INCLUDING N1 AND N2

-- SOME EXPLICIT ASSUMPTIONS THAT MIGHT BE MADE ARE:

@ N1 <= N2 @ @ N1 > 0 @ @ N2 <= N @
I=N1-1 S=0.
```

```
DO I=I+1 # I > N2 # S=S+A(I) OD
(-- IF ONE DIDN'T MAKE THE AFOREMENTIONED ASSUMPTIONS, ONE SHOULD BE
SOMEWHAT MORE CAREFUL IN SOLVING THE LAST PROBLEM --)

-- INITIALIZE LOOP INDEX AND ESCAPE CRITERION
IF N1<=N2 , I=N1-1 IMAX=N2 ; I=N2-1 IMAX=N1 FI
-- CHECK VALIDITY OF LOOP INDEX AND ESCAPE CRITERION
IF I<0 , I=0 FI IF IMAX>N , IMAX=N FI
S=0.  DO I=I+1 # I>IMAX # S=S+A(I) OD

(-- ONE CAN SEE HERE THAT ALTHOUGH THE GENERAL SOLUTION TO THE LAST
PROBLEM REQUIRED A LITTLE THOUGHT, THE PROGRAMMING WAS QUITE EASY --)
```

15. Debugging

Salome doesn't have any facility for debugging, but it does provide a

mechanism for ignoring or not ignoring certain sequences of executable

code at translation time.

One may prefix a '%' to a line of Salome code or one may surround multiple

lines of Salome code with the delimiters '(%' and '%)'.

Salome code which is so delimited may be ignored or not ignored at

translation time depending on whether the '%-Off' or '%-On' statements

have been invoked, respectively. The default is %-Off.

The advantage of this facility with regard to debugging follows.

In the process of debugging, the programmer will have to insert various

write statements, etc., into his program in order to ascertain where the

program is going wrong. This is going to be especially true if the

programmer hasn't taken advantage of the explicit assumption or assertion

statement.

When a Fortran programmer thinks that he has all the bugs in his program licked, he may either (1) remove the debug statements and hope that he won't have to insert them again or (2) make a comment out of each and every debug statement and hope that he won't have to change them all back to executable code.

The Salome programmer may simply insert his debug statements with % delimiting. During debugging, he may activate all or some of these debug statements using the keywords '%-On' and '%-Off'. When a Salome programmer thinks he has all the bugs in his program licked, he may simply eliminate all '%-On' strings from his program and leave all the debug statements in place for possible future use.

Ex.

```
SUB INTERP ( N X Y XI YI ) -- INTERPOLATION ROUTINE
  .
  .
  .
% W ( ' ENTERING INTERP' ) .
@ N > 1 @ -- ASSUME THAT THERE ARE AT LEAST TWO POINTS IN DATA
  .
  .
  .
DX=X(I+1)-X(I)
(% IF DX = 0. , W ( ' DX=0. IN INTERP' ) EXIT /
      DX < 0. , W ( ' DX<0. IN INTERP' ) EXIT FI %)
@ DX > 0. @
D=DY/DX
  .
  .
% W ( ' EXITING INTERP' ) .
RET END
```

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